

## ▼ Motivation

### ▼ Examples

- 4 <-> 9
- Sensor Example
- NYC Taxi Cabs -> Hurricane Sandy vs \$100 tip vs Dropoff in Brazil

### ▼ Core problem: There is no longer one interpretation of the data

#### ▼ Current state of the art:

##### ▼ Design a schema to account for uncertainty

- **Problem:** Now users need to be explicitly aware of uncertainty
- **Problem:** Slow, upfront work

##### ▼ Settle on one interpretation that works for your use case

- **Problem:** If the interpretation you pick is wrong, you get errors
- **Problem:** The data could be wrong if used for a different use case
- **Problem:** Slow, upfront work

##### ▼ NULL values

- **Problem:** Hides uncertain values

##### ▼ **Problem:** Null value semantics are awful

- Any arithmetic with a null value (e.g., NULL + 1) evaluates to NULL
- Any comparison with null values (e.g., NULL >= 3) evaluates to UNKNOWN
- 3-Valued Boolean Logic: TRUE, UNKNOWN, FALSE

##### ▼ SQL WHERE returns only TRUE values (UNKNOWN and FALSE are dropped)

- **Problem:** It's possible for `SELECT * FROM R WHERE (X > 3) AND (X <= 3)` to return an empty result on a non-empty R

#### ▼ Improved Solution: API for Uncertain/Probabilistic Queries

##### ▼ Query for 'certain' answers

- **Problem:** Uncertain answers may still be useful

##### ▼ Query for the best interpretation

- **Problem:** How do you define "best"?

##### ▼ Query for all possible interpretations

- **Problem:** Hides correlations/anticorrelations

##### ▼ Probabilistic queries as above, but also compute...

- ... marginal probabilities of answers
- ... expectations/variances/other statistical measures of answers
- ... rank of each possible answer (when this makes sense)

## ▼ Possible Worlds Semantics

### ▼ Each interpretation defines one world

- ▼ An uncertain database is actually a set of databases, each representing one interpretation or "possible world"
  - For now, all of these databases share the same schema.
- ▼ How do we define query semantics for a set of possible worlds:
  - Queries should return a set of "possible answers"
- ▼ Naive idea: Run the query independently in each possible world
  - **Problem:** Inefficient. Can be lots of possible worlds.
  - **Problem:** Could be impossible. Can be an infinite number of possible worlds
  - **But...** This still defines a self-consistent set of rules for evaluating queries on uncertain data

## ▼ Representation Requirements

- ▼ Closed
  - There exists a  $Q'$  such that  $Q'(\text{Rep}(D)) == \text{Rep}(Q(D))$
- ▼ Meaningful
  - The representation has to be useful... although for what depends on the application
- ▼ ... or better still Bijective
  - Ideally, it would be nice to be able to reconstruct all possible worlds from the representation.

## ▼ Factorization attempts

- ▼ Three types of uncorrelated uncertainty:
  - Row-level: A row is present precisely half of all possible worlds --- and other than the row, everything else is identical between the two halves
  - Attribute-level: There are  $N$  copies of all worlds where a row is present, differing only in a single attribute which takes  $N$  distinct values ---  $N$  may be infinity
  - Open-world: There are an infinite number of worlds with an unbounded number of rows in them, and we have rules for generating more rows

## ▼ Adding correlations

- Create an integer "world-id"
- ▼ Define a function that maps the world-id to a concrete database (or relation) instance
  - ... so how do we define these functions?

# ▼ V-Tables

## ▼ Null Value Semantics on Steroids

- 'Label' each Null. i.e., Nulls become Variables
- ▼ A V-table is effectively a Function:
  - A possible world is defined by a mapping from labels to nulls
  - Externally provided ruleset defines what's allowed to be in a labeled null

## ▼ Proving Closure for V-Tables

- Exercise for the reader
- ▼ Works for  $\pi$ ,  $\times$ ,  $\cup$ , but not  $\sigma$

- ... because there's no way to represent a row that "might" be in the result set
- ▼ Works under both set and bag semantics
  - ... although the representation may have some duplicate rows that need to be removed

## ▼ C-Tables

### ▼ V-Tables with an additional "Condition" column

- ▼ Each table gets an added column containing a boolean expression that may reference label symbols
  - When evaluating the V-Table as a Function, plug label values into the boolean expression
  - Boolean expressions that evaluate to false are not present in that specific possible world.

### ▼ Proving Closure for C-Tables

- Also an exercise for the reader
- ▼ Works for  $\pi$ ,  $x$ ,  $U$ ,  $\sigma$ ,  $\delta$  but not generalized  $\pi$  or  $\gamma$ 
  - ▼ ... well, not entirely true. It works if  $\pi$  and  $\gamma$  are allowed to create new variable symbols and constrain their values based on the values of other symbols
    - ... which means  $\pi$  and  $\gamma$  effectively have side effects
  - Works for both bag and set representations, although as before there may be duplicates

### ▼ Simplified C-Tables (U-Relations)

- Remove Support for Labeled Nulls
- ▼ Create one row for each possible value and add to the condition column `AND [label] = [value]`
  - ... only works if you have a finite, discrete set of possible values
- ▼ Worldset-Decompositions
  - Store the U-Relation column-store style.

### ▼ Generalized C-Tables

- ▼ Allow the creation of new variable symbols defined by formulas
  - e.g.,  $\{ X + 2*Y \}$
- ▼ Closed over SPJUA+Distinct
  - ... although for aggregates/distinct the representation can get very very very large

## ▼ Weaker Models

### ▼ OR-SET encoding

- Label tuples that are not in at least one possible world with a ? (this alone is generally called Tuple-independent)
- Use sets of allowable values instead of attributes
- Can not capture correlations

### ▼ X-Tuples

- Group tuples into sets of mutually exclusive possibilities (can be combined w/ OR-SET)

## ▼ Queries on C-Tables

## ▼ Basic query types

### ▼ Certain Answers

- Answers in \*all\* possible worlds

### ▼ Possible Answers

- Answers in \*any\* possible world

## ▼ Limitations

- Expensive to compute either of these
- Possible produces too much, while certain produces too little.

## ▼ Tradeoff Points

- Best Guess (Maximal Prior) - Pick a (most likely) world and evaluate the query in it
- Maximal Posterior - Use probabilities (discussed next class) to pick result rows exceeding a given threshold probability.
- Sampling - Pick a set of possible worlds at random and evaluate the query in each of those (more discussed soon)